Random Number Generation on a Decimal Computer

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I propose a rapid algorithm to generate uniformly distributed random numbers that is suitable for a decimal computer, specifically the IBM 1401 computer. I employ simple statistical methods to demonstrate the viability of this algorithm. Then I show that it is easy to convert the uniformly distributed random numbers to normally distributed random numbers., i.e., numbers that fit a bell curve.

1. Introduction

The IBM 1401 was an early 1960s era transistor-based small-business computer. It supported a maximum of 16 KB of core memory, and it had an 87 KHz machine cycle – extremely miniscule and painfully slow compared to today's computers. It had a decimal-based architecture and could store a single digit 0 - 9 as an 8-bit character. Each digit took 6 bits of the character, plus there was a parity bit and a "word mark" bit. Numbers were arbitrarily long; a number was addressed by its rightmost digit and the leftmost digit had its word mark bit set. Today, we would call each character a "byte", each number a "string decimal", and say that the IBM 1401 performed "string decimal arithmetic".

The IBM 1401 had basic add and move machine instructions, but no bitwise shift or "and" or "or" instructions. Multiplication and division were slow operations, and floating-point operations were done by software routines. Therefore, an efficient algorithm designed to run on that machine should rely primarily on adds and moves. Arithmetic operations on numbers and moves of numbers from one memory location to another were delimited by word marks. There were machine instructions to set and clear a word mark as needed at any particular memory address.

2. An algorithm to generate uniformly distributed random values

My random number algorithm relies on rotating 7-digit numbers to the right by 3 digits and by 5 digits. Note that 3, 5, and 7 are prime numbers.

To rotate a 7-digit number to the right by 3 digits on the IBM 1401, store the number in memory at the left end of a 10-character buffer. For example, to rotate the number **1234567**:

1234567xxx

Move the number right 3 characters to the right:

1231234567

Move the rightmost 3 characters of the buffer to the beginning of the buffer:

5671234567

Then the leftmost 7 characters of the buffer contain the rotated number **5671234**. Therefore, each rotation requires two moves.

Rotating a 7-digit number to the right is similar, except that you start with the number at the left end of a 12-character buffer. For example, rotating the number **1234567** yields **3456712**.

Steps of the algorithm:

- 1. Start with two nonzero 7-digit seed values, call them **r1** and **r2**.
- 2. Rotate **r1** to the right by 3 characters and rotate **r2** to the right by 5 characters.
- 3. Add the two rotated numbers to product a 7-digit **sum** (drop any overflowed digits at the left).
- Digits at the left end of sum constitute a random value. If, for example, you want 2-digit random numbers, take the leftmost two digits.
- 5. Set **r1** to the value of **r2**, and set **r2** to the value of **sum**.
- 6. Go back to step 2 and iterate as many times as desired. Each iteration produces one random value.

To the right is an example of the first three iterations to generate 2-digit random values.

```
r1 = 1234567
             # seed
r2 = 8901234
             # seed
ITERATION #1
r1
   = 5671234  # rotated right 3
r2 = 0123489 # rotated right 5
sum = 5794723 \# sum = r1 + r2
*** random value = 57
r1 = 0123489 \# r1 = r2
r2 = 5794723 \# r2 = sum
ITERATION #2
   = 4890123  # rotated right 3
r1
r2 = 9472357 # rotated right 5
sum = 4362480 \# sum = r1 + r2
*** random value = 43
r1 = 9472357 \# r1 = r2
   = 4362480  # r2 = sum
r2
ITERATION #3
r1 = 3579472 # rotated right 3
r2 = 6248043 # rotated right 5
sum = 9827515 \# sum = r1 + r2
*** random value = 98
r1 = 6248043 \# r1 = r2
r2 = 9827515 \# r2 = sum
```

3. Python implementation

I implemented the algorithm in Python and generated a million 2-digit random values. I used Python's standard integer datatype, so I did the rotations with integer multiplication, division, and modulo operations with powers of 10. As mentioned above, on the IBM 1401, only moves are required to do the rotations.

```
rvs = [] # list to contain the random values, initially empty
r1 = 1234 567
                # seed
r2 = 89 \ 0\overline{1}234
               # seed
# Loop to generate a million random values.
for i in range(1 000 000):
    r1 = 10000 * (r1 \$1000) + r1//1000
                                        # rotate right 3 digits
    r2 = 100*(r2%100000) + r2//100000 # rotate right 5 digits
                                # add, and keep the sum to 7 digits
    sum = (r1 + r2) \$10000000
    rvs.append(sum//100000)
                                # append the leftmost 2-digit number to list rvs
    r1 = r2
    r2 = sum
```

```
print(f'The first 20 random values:')
print(f'{rvs[:20]}')
```

```
The first 20 random values:
[57, 43, 98, 31, 47, 30, 81, 92, 80, 25, 65, 4, 31, 30, 48, 60, 40, 4, 16, 5]
```

4. Basic statistical confirmation

We need to confirm that we did indeed generate uniformly distributed random values. First, some basic statistics:

```
import statistics
import numpy as np
print(f' mean = {statistics.mean(rvs):.2f}')
print(f' first quartile = {np.percentile(rvs, 25)}')
print(f'second quartile = {np.percentile(rvs, 50)}')
print(f' third quartile = {np.percentile(rvs, 75)}')
```

```
mean = 49.50
first quartile = 25.0
second quartile = 50.0
third quartile = 74.0
```

It's indeed a good sign that the mean (arithmetic average) is very close to 50, the first and second quartiles hit their marks, and the third quartile is off by only 1.

Here's a frequency bar chart of all the values. They do appear to be uniformly distributed!

```
import matplotlib.pyplot as plt
import seaborn as sns
def show chart(title, x label, y label, x values, y values, size=(10, 5)):
    .....
   Display a bar chart.
    @param title the chart title.
    Oparam x label the label for the x axis
    Oparam y label the label for the y axis
    @param x_values the x values to plot
    @param y values the y values to plot
    @param size the size (width, height) of the chart
    .....
    figure = plt.figure(figsize=size)
    axes = sns.barplot(x_values, y_values, color='green', alpha=0.75)
    axes.set title(title)
    axes.set(xlabel=x label, ylabel=y label)
```



Even though the final result appears valid, let's make sure that the values were generated in random order not in clumps or in a pattern. Here are scatter plots that show the generation order of the first 100 generated values, 100 values near the middle, and 100 values near the end:

```
xs = [list(range(100))]
ys = rvs[0:100]
plt.scatter(xs, ys)
plt.show()
```





The values in these three samples do appear to have been generated in random order.

5. Generate normally distributed random values

It is very simple and efficient to extend the above algorithm to also generate normally distributed random values, values that fit a bell curve.

As the uniformly distributed random values are being generated, add them together in groups of one hundred – the first through 100th values, the 101st through 200th values, etc. Keep all the digits of each sum. Compute the average of each group by dividing its sum by 100. Of course, on the IBM 1401, that "division" is done by not including the last two digits of the sum.

According to the Central Limit Theorem of statistics, these group averages are normally distributed random values.

In the Python implementation, I computed group averages separately at the end using the uniformly distributed random values stored in list **rvs**:



6. IBM 1401 Autocoder implementation

[Added November 20, 2019]

Autocoder program by Michael Albaugh. Assembled and run under ROPE 1401 simulation.



```
job 1401 PRNG FROM RON MAK
* Assemble with command
* autocoder -1 MakRand.lis -o MakRand.obj -bV ctload.ac
         CTL 6111
* INDEX REGISTERS
         ORG 87
         DCW @000@ POSITION IN LP BUFFER
X1
         ORG 92
X2
         DCW @000@
         ORG 97
         DCW @000@
х3
         ORG 333
* WORKING STORAGE
       DCW @000000@
ACC
         DCW @000000@
SUM
*
* USES TWO 7-DIGIT SEEDS
R1 DCW @1234567@
R2
         DCW @8901234@
*
* TEST CODE
         ORG 400
START
         MCW @000@,X1
LLOOP
         в
              MAKRND
         MCW SUM-5,203&X1
         SBR X1,3&X1
         С
              @060@,X1
         BU
             LLOOP
         W
         CS
              332
         CS
         н
* PRNG PROPER
         ORG 500
         SBR MRRTN&3 SAVE RETURN ADDRESS
MAKRND
* ROTATE R1 RIGHT BY THREE CHARACTERS.
         MCW R1-3, ACC MOVE NUM MSDS TO ACC LSDS
         MCW R1
                          MOVE NUM LSDS TO ACC MSDS (CHAIN)
         MCW ACC, R1
* ROTATE R2 RIGHT BY FIVE CHARACTERS
         MCW R2-5,ACC
                         MOVE NUM MSDS TO ACC LSDS
         MCW R2
                          MOVE NUM LSDS ACC MSDS (CHAIN)
         MCW ACC, R2
         MCW R2, SUM
             R1,SUM
         Α
         MZ
              000,SUM-6
                         IGNORE OVERFLOW
         MCW R2,R1
         MCW SUM, R2
MRRTN
         в
              MRRTN
                       RETURN TO CALLER
         END START
```